

# Enabling Secure Machine-to-Machine Interactions in Industrial Internet of Things Ecosystems

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**Abstract**—In the emerging Industrial IoT era, Machine-to-Machine (M2M) communication technology is considered as a key underlying technology for building Industrial IoT environments where devices (e.g., sensors, actuators, gateways) are enabled to exchange information with each other in an autonomous way without human intervention. However, most of the existing M2M protocols that can be also used in the Industrial IoT domain provide security mechanisms based on asymmetric cryptography resulting in high computational cost. As a consequence, the resource-constrained IoT devices are not able to support them appropriately and thus, many security issues arise for the Industrial IoT environment. Therefore, lightweight security mechanisms are required for M2M communications in Industrial IoT in order to reach its full potential. As a step towards this direction, in this paper, we propose a lightweight authentication mechanism, based only on hash and XOR operations, for M2M communications in Industrial IoT environment. The proposed mechanism is characterized by low computational cost, communication and storage overhead, while achieving mutual authentication, session key agreement, device's identity confidentiality, and resistance against the following attacks: replay attack, man-in-the-middle attack, impersonation attack, and modification attack.

**Index Terms**—Industrial IoT, M2M communications, Lightweight Authentication, Security, Sensors.

## I. INTRODUCTION

THE Industrial Internet of Things (IIoT) technology is a key enabler for the next industrial revolution, known as Industry 4.0 [1], [2]. The invention of the steam engine by James Watt in the 18th century caused the first generation of industrial production; the invention of electric power brought about the second industrial revolution in 1870

with the widespread use of electric machines in production lines [3]. The industrial automation and widespread adoption of computers and programmable logic controllers (PLC) in 1970s staged the third industrial revolution. Nowadays, IIoT and Machine-to-Machine (M2M) communications are about to bring the fourth industrial revolution known as Industry 4.0 [4]–[6], where man, machine, and product will be interconnected throughout the

whole supply-chain from the production floor to the managerial level. This will boost the productivity and allow customised and flexible production while benefiting from the economies of scale [7], [8].

However, the transition from the third industrial revolution to Industry 4.0 raises a wide spectrum of new security issues [3], [5], [9]–[11]. Traditional industrial communication systems have been designed for reliable operation in a noisy factory environment, employing mainly hard-wired propriety-based communication technologies to connect sensors, actuators, and controllers as well as other industrial components such as Supervisory Control and Data Acquisition System (SCADA) and Manufacturing Execution System (MES). Nevertheless, with the emergence of IIoT, future factories will increasingly rely on diverse communication technologies including wireless standards to ensure connectivity, interoperability, and remote operation and control of production processes through the Internet. This provides an unprecedented attack surface for the attackers. Unlike computer networks, where an attack normally threatens the information integrity, confidentiality, or availability, attacks against a smart factory can cause physical damage, threaten the human life, render the quality or final products by compromising the production processes, or lead to increased use of resources [9]. Last but not least, unlike the short lifetime of consumer electronics, the lifespan of machines operating on a production floor normally lasts for several decades, and it is not always economically viable to completely replace the legacy equipment with the latest technology. Hence, it is essential to come up with novel solutions that ensure security not only for the leading-edge manufacturing technology but also for the legacy systems.

Since authentication is the cornerstone of providing effective security, a number of authentication schemes have been proposed to ensure security in IoT or M2M applications [12]– [16]; however, they cannot be readily



applied for IIoT because manufacturing machines are naturally limited with computation power and/or communication bandwidth. In this paper, we propose a lightweight authentication scheme to authenticate these resource constrained machineries in order to ensure secure integration of IIoT solutions in the future production systems.

To this end, we consider an IIoT scenario where a machine

(i.e., a smart sensor), equipped with a Secure Element (SE), is authenticated by a network element (i.e., a router) equipped with a Trusted Platform Module (TPM).

The rest of this paper is organized as follows. Section II reviews the related work. Section III describes the proposed authentication mechanism. Section IV presents the security analysis for the proposed mechanism, and Section V discusses its performance evaluation in terms of communication overhead, computational cost, and storage cost. Finally, Section VI concludes.

## II. RELATED WORK

### A. M2M Communication protocols for IIoT

The main focus of communication protocols is to guarantee the delivery, routing and storage of the information without the necessity of implementing different mechanisms in different devices or applications. Based on relevant scientific works, we evaluate the most popular protocols for use in IIoT and M2M communication. As a result, according to [33] and [34] the most discussed and promoted standards for the communication between devices and cloud services are:

- 6LoWPAN: IPV6 over Low power Wireless Personal Area Network

An Internet Protocol version 6 (IPv6) over low-power wireless personal area networks (6LoWPAN) standard has been developing, by the Internet Engineering Task Force (IETF), in order to promote the development of the IoT and exploit the M2M applications. 6LoWPAN enables IP-based M2M devices to connect to the Internet. More precisely, one potential application is to monitor the manufacture process in industry [26], [27], where a number of sensors, actuators, and controllers are connected together to achieve passive monitoring and active control and automation. However, 6LoWPAN has various security challenges and many threats and trust crises are existing along with its development. Sensor nodes are usually distributed in an unprotected environment and messages can be easily eavesdropped in the transmission. A lot of

research works have been proposed to overcome the vulnerabilities of the 6LoWPAN systems. A secure authentication and key establishment scheme (SAKES) has been proposed in [17]. SAKES makes use of two different cryptographic schemes. Firstly, the authors propose the use of a symmetric key scheme to encrypt messages in the authentication phase and then, an asymmetric key scheme based on the elliptic curve cryptography (ECC) is used, in the key establishment phase, to build a session key between the 6LoWPAN devices and the server. More precisely, the session key in the SAKES scheme is calculated based on the Diffie–Hellman (DH) key exchange method. EAKES6Lo scheme proposed by Qiu et.al is a mutual authentication and key establishment scheme for M2M communication in 6LoWPAN Networks [18]. It consists of three phases: pre-deployment phase, AKE phase, and handover phase. Their security analysis shows that the proposed scheme can be resistant against replay attacks, Man-in-the-Middle attacks, impersonation attacks, Sybil attacks, and compromised attacks. Moreover, it can benefit from low computational overhead and transmission overhead.

- MQTT: Message Queue Telemetry Transport Message Queue Telemetry Transport, is an OASIS standardized protocol [19]. It is designed to be lightweight, flexible and simple to implement. MQTT uses different routing mechanisms, such as: one-to-one; one-to-many or many-to-many, making possible the connection for IoT and M2M to connected devices/applications. MQTT is a publish/subscribe messaging transport protocol on the top of TCP/IP protocol consisting of three components (subscriber, publisher, and broker). In terms of Quality of Service (QoS), it supports three levels: (i) QoS-0: The message will be delivered once, with no confirmation, (ii) QoS-1: The message will be delivered at least once, with confirmation required, (iii) QoS-2: The message will be delivered exactly once by using a handshake. However, although MQTT has been deployed for IoT, it has limited security features addressing IoT security issues. In particular, MQTT uses user-name/password authentication and SSL/TLS for secure data communication [20]. Hence, further efforts are required so that MQTT can address effectively and efficiently security issues for IoT. In [11], the authors have envisaged the use of SSL/TLS with certificates and session key management to enhance security for MQTT. In particular, they have proposed a Secure MQTT (SMQTT). Their proposed solution is based on lightweight Attribute Based Encryption (ABE) [22], [23] over elliptic curves [35]. They have used ABE

because of its inherent design supporting broadcast encryption that is suitable for IoT applications. ABE includes two types: (i) Cipher-text Policy based ABE (CP-ABE), and (ii) Key Policy based ABE (KPABE). Authors' analysis and performance evaluation show that MQTT is efficient, robust, and scalable.

- **AMQP:** Advanced Message Queuing Protocol Advanced Message Queuing Protocol, is a binary application layer protocol standardized from the ISO/IEC 19464. It is a message centric protocol on top of TCP/IP, which provides publish-subscribe and point-to-point communication. AMQP supports message-oriented communication via message-delivery guarantees including: (i) at-most-once, when each message is delivered once or never, (ii) at-least- one, when each message is delivered and (iii) exactly-one, when the message will certainly delivered only once [24]. AMQP provides different features, including routing and storing messages within the broker using message queues. In terms of security, it supports SASL authentication and TLS for secure data communication [25].
- **CoAP:** Constrained Application Protocol Constrained Application Protocol, is a web transfer protocol which supports unicast and multicast requests for use in constrained devices and networks. It is based on a request-response architecture between endpoints. CoAP clients after sending the requests using an URI, can receive as a response GET, PUT, POST and DELETE resources from the server [26]. The messages are

exchanged over UDP between endpoints and also it supports the use of unicast and multicast requests. In terms of QoS, CoAP supports two levels: (i) 'confirmable' when no packet is lost and the receiver respond with an ACK; and (ii) 'non-confirmable' when the message do not require an ACK. CoAP provides security via the Datagram Transport Layer Security (DTLS) which is a secure protocol for network traffic to support handling packet loss, reordering of messages and message size [27]. However, DTLS requires numerous message exchanges to establish a secure session and thus it is characterized by high communication cost. Therefore, CoAP suffers from this challenge of DTLS. In order to overcome this issue, a lightweight secure CoAP for the Internet of

Things (Lithe) scheme has been proposed by Raza et.al [28]. Lithe scheme shows that DTLS can be compressed and its overhead is significantly reduced. The proposed scheme is implemented in Contiki [29] and their evaluation shows significant gains in terms of packet size, energy consumption, processing time, and network-wide response times when compressed DTLS is enabled.

- **XMPP:** Extensible Messaging and Presence Protocol Extensible Messaging and Presence Protocol, is a TCP communication protocol based on Extensible Markup Language (XML) used for real-time messaging, online presence and request-response services [30]. Clients communicate via a distributed network and do not rely on a central broker. XMPP supports publish/subscribe model and provides

TABLE I OVERVIEW OF POSSIBLE M2M PROTOCOLS FOR IIOT

Protocol	Feature	Layer	TCP/UDP	Security	Related work
6LOWPAN	To map services required by the IPv6 over Low power WPANs to maintain an IPv6 network	Network	TCP	SSL	[17], [18]
MQTT	To utilize the publish/subscribe pattern to provide transition flexibility and simplicity of implementation	Application	TCP	SSL	[19]–[23]
AMQP	To provide publish-subscribe and point-to-point communication	Application	TCP	SSL	[24], [25]
CoAP	To connect resource-constrained devices in a secure and reliable way	Application	UDP	DTLS	[26]–[29]
XMPP	To transfer instant messaging (IM) standard that is used for multi-party chatting, voice and video calling and telepresence	Application	TCP	SSL	[30], [31]
DSS	To enable scalable, real-time, dependable, high-performance and interoperable data exchanges using a publish–subscribe pattern	Application	TCP/UDP	SSL	[32], [33]

security such as authentication via SASL and secure



$ID_i$	The identity of smart sensor $i$
$AID_i$	The alias of entity $i$
$f_i$	Function generation
$SK_i$	The shared key between smart sensor $i$ and a router
$R_i$	A random number generated by a Pseudorandom Number Generator (PRNG)
$h(\cdot)$	A one-way hash function
	A concatenation operator
	XOR operation

A. Registration

Each smart sensor needs to perform the registration procedure with the AS through a secure channel. The AS generates the secure pre-shared key set  $PSK_i, i = 1, \dots, n$  and sends each  $PSK_i$  to one of the routers. The registration procedure consists of the following steps, as it is shown in Figure 1.

- 1) Smart sensor  $\rightarrow$  AS: The smart sensor transmits its unique identity number  $ID_i$  to the AS over a secure channel.
- 2) Upon receiving the smart sensor's ID, the AS calculates the following three secret authentication parameters for the sensor:  $f_{ii} = h(ID_i \text{ k } x)$ ,  $f_{2i} = h(f_{ii})$ , and  $f_{3i} = PSK \oplus f_{ii}$ . The objective of  $f_{ii}, f_{2i}$  is to build the relation between sensor's ID and AS.
- 3) AS  $\rightarrow$  Smart sensor: The AS sends the following parameters in the smart sensor via a secure channel:  $f_{2i}, f_{3i}$ . These parameters are stored in the SE of the sensor.

B. Authentication

After the registration phase, each sensor is able to authenticate to a router. It is worthwhile to mention that, during the authentication procedure, the sensor never uses its real identity to authenticate to a router. Hence, the smart sensor's ID cannot be eavesdropped by a malicious entity. The authentication procedure consists of the following steps, as it is shown in Figure 2.

- 1) The smart sensor generates a random number  $R_i$  and stores it in the SE of the sensor. Then, it computes the parameter  $M_1$  as follows:  $M_1 = h(f_{2i}) \oplus R_i$ . Afterwards, the sensor computes its alias as  $AID_i = h(R_i) \oplus ID_i$  and computes the parameter  $M_2$  as follows:  $M_2 = h(R_i \text{ k } M_1 \text{ k } AID_i)$ .
- 2) Smart sensor  $\rightarrow$  router: The smart sensor sends to the router an authentication request (i.e., Message 3) including  $(M_1, M_2, f_{3i}, AID_i)$ .
- 3) Upon receiving the authentication request, the router performs the following. Firstly, the router retrieves  $f_{ii}$  by using the pre-shared key  $PSK$  (i.e.,  $f_{ii} = f_{3i} \oplus PSK$ ). Then, the router obtains  $R_i$  and  $ID_i$  via  $R_i = M_1 \oplus h(f_{2i})$  and

$ID_i = AID_i \oplus h(R_i)$ , respectively. Then, router checks whether  $h(R_i \text{ k } M_1 \text{ k } AID_i)$  is equal to  $M_2$ . The authentication request is rejected if  $h(R_i \text{ k } M_1 \text{ k } AID_i)$  and  $M_2$  do not match. Next, the router generates a random number  $R_2$  that is stored in the TPM of

$$M_1, \text{ and } M_2'$$

the router.

Then, it computes  $AID_j, M_1$  as follows:  $AID_j = R_2 \oplus h(ID_i), M_1 = f_{ii} \oplus h(ID_i),$

and  $M_2 = h(M_1 \text{ k } AID_j \text{ k } R_2)$ . Finally, the router calculates the session key  $SK_{ij}$  as  $SK_{ij} = h(R_i \text{ k } R_2)$ .

- 4) Router  $\rightarrow$  Smart sensor: The router sends back to the sensor the authentication response (Message 4) including  $M_1, M_2$  and  $AID_j$ .

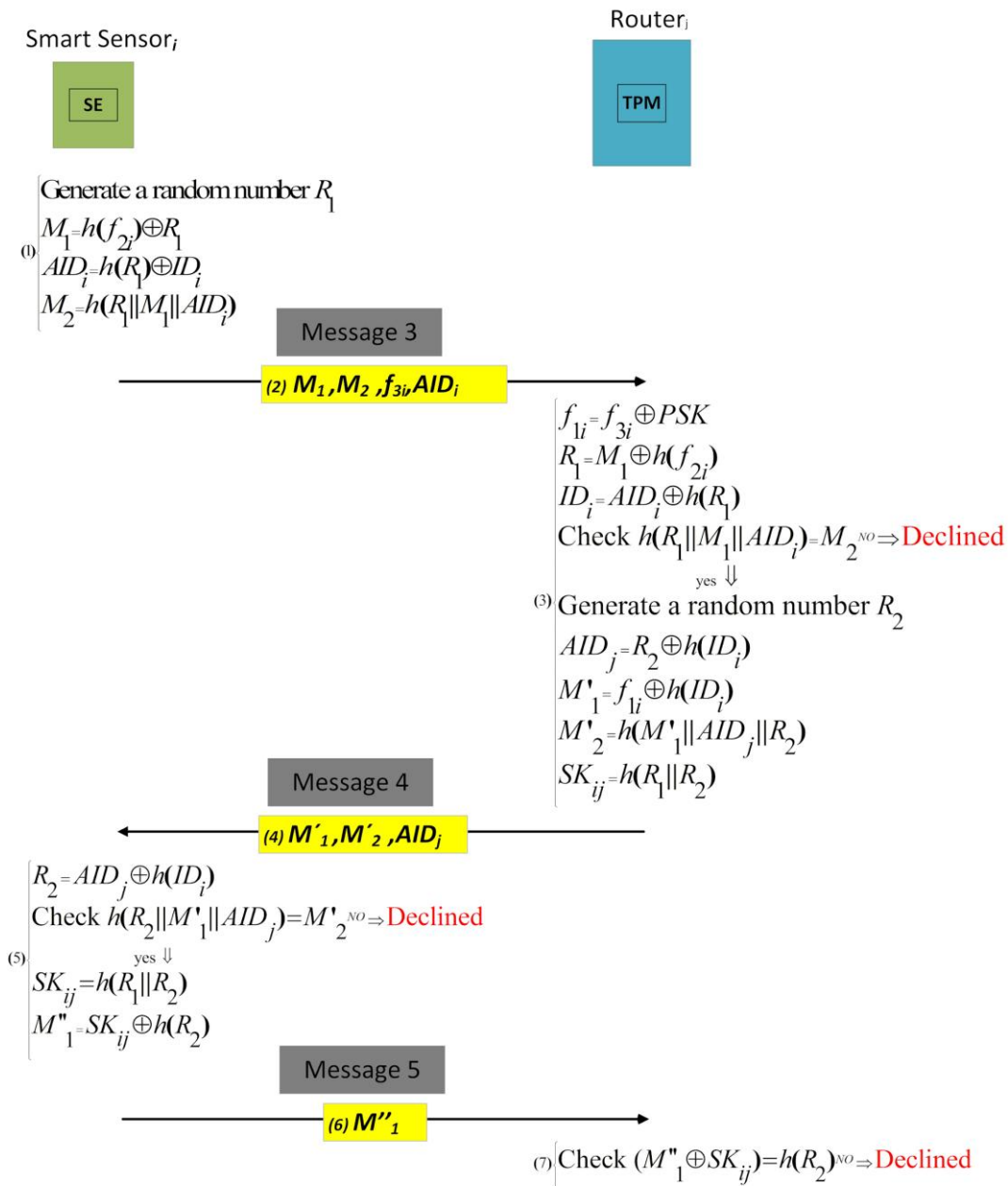


Fig. 2. Authentication procedure.

- The smart sensor retrieves  $R_2$  by computing  $AID_j \oplus h(ID_i)$  and checks if  $h(R_2 || M'_1 || AID_j)$  and  $M'_2$  are equal. If they are equal, the router calculates the session key  $SK_{ij}$  as  $SK_{ij} = h(R_1 || R_2)$ . Finally, the smart sensor calculates  $M''_1$  as  $SK_{ij} \oplus h(R_2)$ .
- Smart sensor  $\rightarrow$  Router: The smart sensor sends Message 5 including  $M''_1$  to the router.

- Upon receiving Message 5, the router uses its session key  $SK_{ij}$  calculated in step 3 in order to retrieve  $h(R_2)$ .

Then, the router calculates  $SK_{ij} \oplus M''_1$  and compares it with  $h(R_2)$ . If they are equal, it means that the smart sensor holds the legitimate session key.

#### IV. SECURITY ANALYSIS

In this Section, we provide the security analysis of the proposed authentication mechanism. We have adopted the security analysis approach followed in [39]–[41] and thus, we have the following:

*Proposition 1. Smart sensor's identity confidentiality is provided by the proposed mechanism.*

*Proof.* In the proposed mechanism, the confidentiality of the smart sensor's identity is based on a hash value of a random number  $R_i$  and a XOR function (i.e.,  $AID_i = h(R_i) \oplus ID_i$ ). Therefore, the adversary cannot derive the identity  $ID_i$  of the smart sensor without knowing the random number  $R_i$ .  $\square$

**Proposition 2.** *The proposed mechanism provides entity mutual authentication.*

*Proof.* In the authentication phase, mutual authentication between the smart sensor and the router can be achieved based on the received Message 3 and Message 4. Upon receipt of  $M_1, M_2, f_{3i}$ , and  $AID_i$ , the router checks whether  $M_2$  is equal to  $h(R_i \parallel M_1 \parallel AID_i)$ . The smart sensor is considered authenticated if the equality holds. The same process takes place in authenticating the router when the smart sensor

receives  $M_1, M_2$  and  $AID_j$ . The smart sensor computes  $h(R_2 \parallel M_1' \parallel AID_j)$  and checks whether this value is equal to  $M_2$ . If they are equal, the router is also considered as authenticated. Moreover, if the adversary aims to forge a valid smart sensor/router, he/she needs to generate valid messages. However, the adversary cannot generate the valid messages because he has no information about the random numbers (i.e.,  $R_1$  and  $R_2$ ).  $\square$

**Proposition 3.** *The proposed mechanism is resistant to replay attack.*

*Proof.* As it is described in Subsection III-B, we assume that a legitimate smart sensor has sent Message 3 (i.e.,  $M_1, M_2, f_{3i}$ , and  $AID_i$ ) to the router. If an adversary tries to impersonate the legitimate smart sensor by replaying Message 3, the router will reject the authentication request because the alias  $AID_i$  of the smart sensor is calculated based on a hash value of a random number  $R_i$  which is only known to the legitimate sensor.  $\square$

**Proposition 4.** *The proposed mechanism is resistant to man-in-the-middle attack.*

*Proof.* By obtaining the smart sensor's identity (i.e.,  $ID_i$ ), an adversary is not able to launch a man-in-the-middle attack and computes a session key  $SK_{ij}$ , because he/she is not able to obtain the secret key  $x$  that is only known to the AS.  $x$  is securely stored in the authentication server and is never transmitted to any other entity. Meanwhile, the adversary cannot pretend that is a trustful router to authenticate other smart sensors since he/she does not have the pre-shared secret key  $PSK$ .  $\square$

**Proposition 5.** *The proposed mechanism is resistant to impersonation attack.*

*Proof.* We consider that a smart sensor (e.g.,  $smartsensor_k$ ) intends to impersonate another smart sensor (e.g.,  $smartsensor_i$ ). However,  $smartsensor_k$  cannot obtain the session key  $SK_{ij}$  generated by the  $smartsensor_i$  and the router because it has no information about the random number of  $smartsensor_i$ .  $\square$

**Proposition 6.** *The proposed mechanism is resistant to modification attack.*

*Proof.* The one-way hash function  $h()$  guarantees that information cannot be modified without being detected. If an adversary transmits a modified message to the router, the router will detect it by checking the hash values.  $\square$

TABLE III

SETTING OF PARAMETERS

Parameters	Value (bits)
$ID_i$	128
$AID_i$	128
Random number	128
Hash value	128
$f_{1i}$	128
$f_{2i}$	128
$f_{3i}$	128

Moreover, the proposed authentication mechanism can achieve the following security objectives:

1) No clock synchronization:

In the proposed mechanism, it is not required to synchronize the clock of the devices (e.g., smart sensor, router, and authentication server) because the messages which are exchanged between the devices are similar to the messages exchanged in the nonce-based authentication mechanism [38] which does not rely on timestamps.

2) Fast error detection:

In the authentication procedure, the router will detect an error immediately if an adversary uses the wrong sensor ID. This means that a legitimate smart sensor will have been identified by the router before the calculation of session key.

3) Independent session key:

If the session key  $SK_{ij}$  is compromised by an adversary, the smart sensor and the router can generate a new session key. This is because the generation of the session key  $SK_{ij}$  in the proposed authentication mechanism is based on a hash function and random numbers and thus, it is independent to the previous session key.

V. PERFORMANCE EVALUATION

In this Section, we evaluate the performance of the proposed authentication mechanism in terms of communication overhead, computational cost, and storage overhead.

A. Communication overhead

To analyze the communication overhead, we assume that the five messages (i.e., *Message 1*, *Message 2*, *Message 3*, *Message 4*, and *Message 5*) are transmitted during the registration procedure and the authentication procedure. More precisely, *Message 1* and *Message 2* are transmitted in the registration procedure and *Message 3*, *Message 4*, and *Message 5* are transmitted in the authentication procedure. *Message 1* includes the sensor's identity  $ID_i$  and *Message 2* includes  $f_{2i}$  and  $f_{3i}$ . In addition, *Message 3* includes  $M_1, M_2$ , and  $AID_i$  which are calculated as follows:

- $|M_1| = |h(f_{2i}) \oplus R_1|$
- $|M_2| = |h(R_1) \parallel M_1 \parallel AID_i|$
- $|AID_i| = |h(R_1) \oplus ID_i|$

Moreover, the parameters  $M_1, M_2$ , and  $AID_j$  in *Message 4*, are calculated as follows:

- $|M_1'| = |h(ID_i) \oplus f_{1i}|$

- $|M_2'| = |h(R_2 \parallel M_1' \parallel AID_j)|$
- $|AID_j| = |h(ID_i) \oplus R_2|$

Finally, the parameter  $M_1$  in *Message 5*, is calculated as follows:

- $|M_1''| = |SK_{ij} \oplus h(R_2)|$

Table III contains the setting of parameters that we have assumed for evaluating the communication overhead of the proposed mechanism. Therefore, the overall bandwidth overhead of the proposed mechanism is calculated as follows:

$$bw = \sum_{i=1}^5 Message\ i \tag{1}$$

- *Message 1* =  $|ID_i| = 128$  bits
- *Message 2* =  $|f_{2i}| + |f_{3i}| = 256$  bits

- *Message 3* =  $|M_1| + |M_2| + |f_{3i}| + |AID_i| = 768$  bits
- *Message 4* =  $|M_1'| + |M_2'| + |AID_j| = 640$  bits
- *Message 5* =  $|M_1''| = 128$  bits

B. Computational cost

To calculate the computational cost of the proposed mechanism, we have considered the following notations:  $C_h$  denotes the cost of one-way hash function;  $C_{XOR}$  denotes the cost of XOR operation; and  $C_{ran}$  denotes the cost of generating a random number. Based on the three components (i.e, smart sensor, router, and authentication server) which are used in the proposed mechanism, the computational cost of each component is presented as follows:

1) Smart Sensor

The smart sensor performs computations only in the authentication procedure. Thus, the total computational cost of this procedure is  $4 * C_{XOR} + 7 * C_h + C_{ran}$ .

2) Router

The router performs computations only in the authentication procedures. Therefore, it needs to perform

$$6 * C_{XOR} + 7 * C_h + C_{ran} \text{ operations.}$$

3) Authentication Server

The authentication server performs  $C_{XOR} + 2 * C_h$  operations to allow a smart sensor to be registered.

The proposed mechanism's computational cost is illustrated in Table IV. Due to the fact that the proposed authentication mechanism is based only on XOR operation and hash operation is efficient in terms of the computational cost.

C. Storage overhead

In the proposed mechanism, only a few parameters are needed to be stored by the smart sensor, the router, and the AS. The smart sensor stores its unique identity number  $ID_i$  and the following parameters:  $f_{2i}, f_{3i}, R_1, R_2, AID_i, AID_j$ ,

$M_1, M_2, M_1', M_2', M_1'',$  and  $SK_{ij}$ . In the other side, the router stores the pre-shared key  $PSK$  and the following parameters:

TABLE IV COMPUTATIONAL COST OF OUR PROPOSED MECHANISM

	Smart Sensor	Router	Authentication Server
Registration	—	—	$C_{XOR} + 2 * C_h$
Authentication	$4 * C_{XOR} + 7 * C_h + C_{ran}$	$6 * C_{XOR} + 7 * C_h + C_{ran}$	—

$f_{1i}, f_{3i}, R_1, R_2, ID_i, AID_i, AID_j, M_1, M_2, M_1, M_2, M_1$ , and  $SK_{ij}$ .

Moreover, the AS stores  $f_{1i}, f_{2i}, f_{3i}$ , and  $PSK$ . Following the setting of parameters presented in Table III, the proposed mechanism's storage cost is illustrated in Table V.

TABLE V  
STORAGE COST OF OUR PROPOSED MECHANISM

	Smart Sensor	Route r	Authentication Server
$ID_i$	X	X	—
$AID_i$	X	X	—
$AID_j$	X	X	—
$R_1$	X	X	—
$R_2$	X	X	—
$M_1$	X	X	—
$M_2$	X	X	—
$M_1^0$	X	X	—
$M_2^0$	X	X	—
$M_1^{00}$	X	X	—
$SK_{ij}$	X	X	—
$f_{1i}$	—	X	X
$f_{2i}$	X	—	X
$f_{3i}$	X	X	X
$PSK$	—	X	X
Total cost (bits)	2176	2304	512

VI. CONCLUSION

In this paper, we have proposed a lightweight authentication mechanism, based only on hash and XOR operations, for M2M communications in Industrial IoT environment. The proposed mechanism is characterized by low computational cost, communication and storage overhead, while achieving mutual authentication, session key agreement, device's identity confidentiality, and resistance against the following attacks: replay attack, man-in-the-middle attack, impersonation attack, and modification attack. As future work, we plan to extend the proposed authentication mechanism to provide lightweight mutual authentication between sensors in Industrial IoT environment.

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